

6 Application: Combinatorial Systems

We now turn our attention from the constraining power of the GR operator to its ability to *generate* a language through its second operand. If the first operand of the GR operator is the universal language, the second operand specifies strings that remain. In comparison, if a logical formula ϕ_1 is a tautology, the truth value of the material implication $\phi_1 \rightarrow \phi_2$ depends completely on the right-hand side ϕ_2 . Accordingly, any language $X \subseteq \Pi^*$ can be passed through the GR operation unchanged as follows: $X = [\Pi^* \xrightarrow{\Pi,0,M} X]$.

6.1 String Coverings with a Lexicon

The things get very interesting when we modify the first operand by adding into it a diamond that occurs before an *arbitrary* character position. This changes a lot: a string in the second operand would now be passed through the GR operation only if its *every* character is disjunctively preceded by a hidden diamond:

$$X = [\nu'_{1,1}(\Pi^*) \xrightarrow{\Pi,1,M} \nu'_{*,1}(X)] \text{ where } \nu'_{1,j}(W) = d_{\{\diamond_1, \dots, \diamond_j\}}^{-1}(W) \cap (\epsilon \cup (\Pi^* M) \Pi^+) \quad (31)$$

$$\text{and } \nu'_{*,j}(W) = d_{\{\diamond_1, \dots, \diamond_j\}}^{-1}(W) \cap (\epsilon \cup (\Pi^* M)^* \Pi^+).$$

The same effect is captured by adding two diamonds that surround each character and empty string on both operand languages:

$$X = [\Pi^* \nu_{2,1}(\epsilon \cup \Pi) \Pi^* \xrightarrow{\Pi,2,M} \nu_{*,1}(X)] \text{ where } \nu_{2,j}(W) = d_{\{\diamond_1, \dots, \diamond_j\}}^{-1}(W) \cap (\Pi^* \diamond_j \Pi^* \diamond_j \Pi^*)$$

$$\text{and } \nu_{*,j}(W) = d_{\{\diamond_1, \dots, \diamond_j\}}^{-1}(W). \quad (32)$$

Free Contexts We can now elaborate the right-hand side of the GR and add there left and right contexts $\# \Pi^* _ \Pi^* \#$:

$$X' = [\nu'_{1,1}(\Pi^+) \xrightarrow{\Pi,1,M} \Pi^* \nu'_{*,1}(X) \Pi^*] = [\Pi^* \nu_{2,1}(\Pi) \Pi^* \xrightarrow{\Pi,2,M} \Pi^* \nu_{*,1}(X) \Pi^*]. \quad (33)$$

Every string in the result X' will now have a *string covering* that consists of possibly overlapping factors X . For example, if $X = \{\mathbf{autom}, \mathbf{mate}, \mathbf{eria}\}$ then X' contains such strings as $\mathbf{automate}$, $\mathbf{materia}$ and $\mathbf{automateriaautom}$. Accordingly, we have defined a simple combinatorial system. If we want, we can concatenate a unique *sentinel* symbol $\sigma \in \Pi$ to the end of the lexicon $X \subseteq (\Pi \setminus \sigma)^*$ in order to generate $X' = (X\sigma)^*$ *i.e.* a set of strings covered with non-overlapping factors taken from set $X\sigma$.

Problem 2. Can coverings be used to describe allomorph selection, nonconcatenative morphotactics, interdigitation and multi-component rewriting?

Problem 3. Can $\nu'_{*,j}$ or $\nu_{*,j}$ be used on the left hand side of an extended GR operator that would still preserve star-freeness? How the change interacts with star-freeness, automata size and applications?